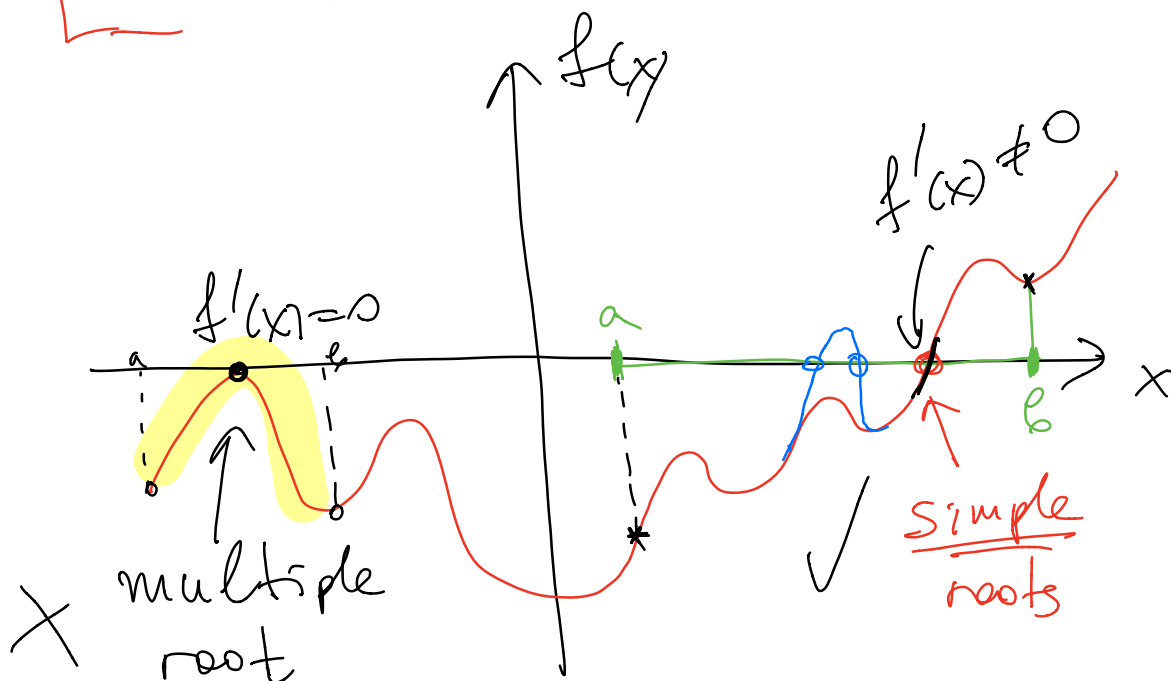


# Nonlinear equations (in one variable)

$$x = \sqrt{c} \quad (\Leftrightarrow) \quad x^2 - c = 0$$

$$\cos x + x^2 - 7 = 0$$

Solve  $f(x) = 0 \quad x \in [a, b]$   
Find at least one solution  
if one exists



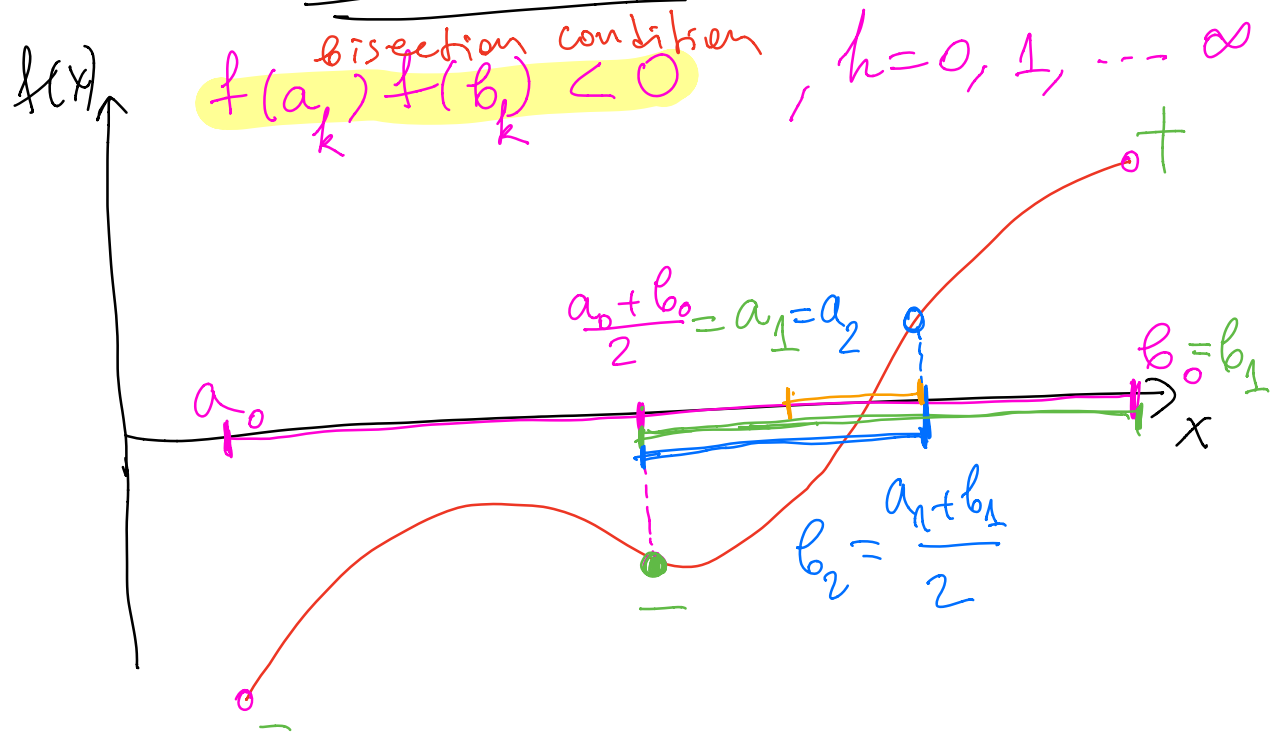
$f(x)$  is continuous on  $[a, b]$

$\exists f$   $f(a) \cdot f(b) < 0$

$\Rightarrow \exists x \in [a, b]$  s.t.  $f(x) = 0$

(Q) Why don't I write  $f(a) f(b) \leq 0$  ?

### Bisection method



## Algorithm

Input:  $[a, b]$  s.t.  $f(a)f(b) < 0$   
 $k_{\max} \in \mathbb{Z}^+$

Output:  $\begin{cases} \tilde{x} \text{ s.t. } f(\tilde{x}) \approx 0 \\ [a, b] \text{ s.t. true root is in } [a, b] \\ x \in [a, b], f(x) = 0 \end{cases}$

For  $k=0, 1, 2, \dots, k_{\max}$

$$x_k = \frac{a_k + b_k}{2}$$

$$f_k = f(x_k)$$

(actually, we only need  
Sign of  $f_k$ )

If  $f_k \cdot f(a_k) < 0$  then  
left half  $[a_k, x_k]$  is a bisection interval,  
 $x \in [a_k, x_k]$

$$a_{k+1} = a_k ; b_{k+1} = x_k$$

else  $\rightarrow$  [if  $f_k \cdot f(b_k) < 0$ ]  
(keep right half)

$$b_{k+1} = b_k ; a_{k+1} = x_k$$

end if

else if  $f_k = 0$  then  
return  $x_k$

End for

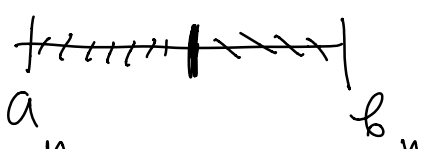
Output :  $x \approx x_{k_{\max}} = \frac{a_{k_{\max}} + b_{k_{\max}}}{2}$

and  $x \in [a_{k_{\max}}, b_{k_{\max}}]$

We know that  $[n = k_{\max}]$

$$|x - x_n| \leq \frac{b_n - a_n}{2} = \frac{b_0 - a_0}{2 \cdot 2^n}$$

$x_n = \frac{a_n + b_n}{2}$



Absolute error

$$e_n = |x - x_n| \leq \frac{b - a}{2^{n+1}}$$

error estimate

Given an error tolerance

$$\epsilon \text{ s.t. } |x - x_n| < \epsilon$$

$$\epsilon \approx \frac{b - a}{2^{n+1}}$$

$$2^{n+1} > \frac{b-a}{\epsilon}$$

$$n+1 = \log_2 \left( \frac{b-a}{\epsilon} \right)$$

How many  
times FO  
will evaluate f(x)

$$n = \left\lceil \log_2 \left( \frac{b-a}{\epsilon} \right) \right\rceil$$

round  
up

/ceil

in Matlab

round  
down

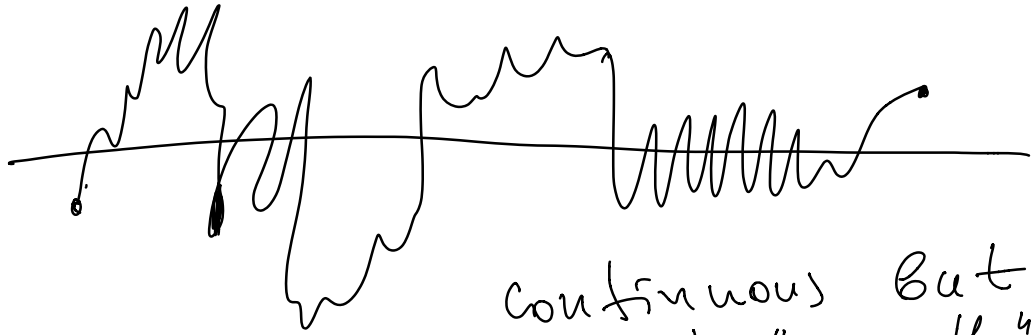
→  $\lfloor \quad \rfloor$  / floor

In Matlab :

$$n_{\text{est}} = \text{ceil}(\log_2((b-a)/\epsilon))$$

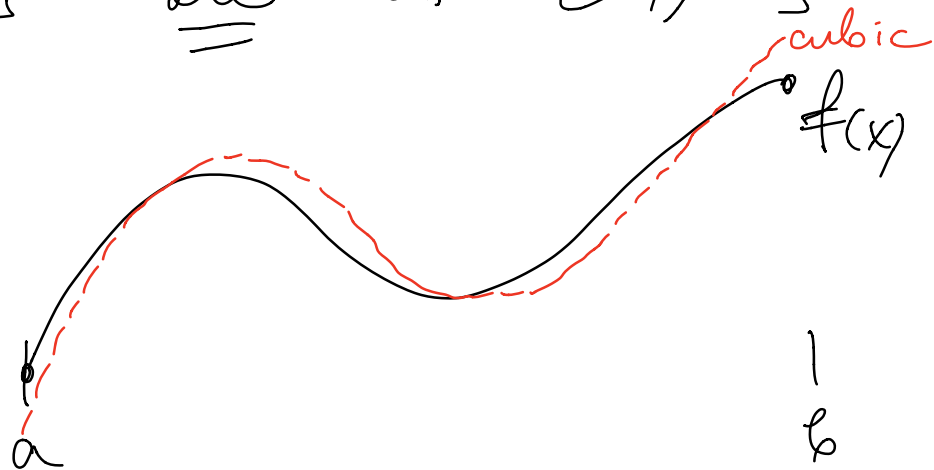
Our job is done

Can we do better?



continuous but  
not "smooth"

$f(x)$  is "smooth" <sup>on  $[a, b]$</sup>  if it  
can be approximated "well"  
by a polynomial of low  
degree (linear, quadratic, cubic)  
over all of  $[a, b]$



Another (related) is that  
 $f(x)$  is smooth if it  
is sufficiently differentiable  
 $e^{1/x}$  is a bad example

---

Idea:  $p(x) \approx f(x)$   
polynomial on  $[a, b]$

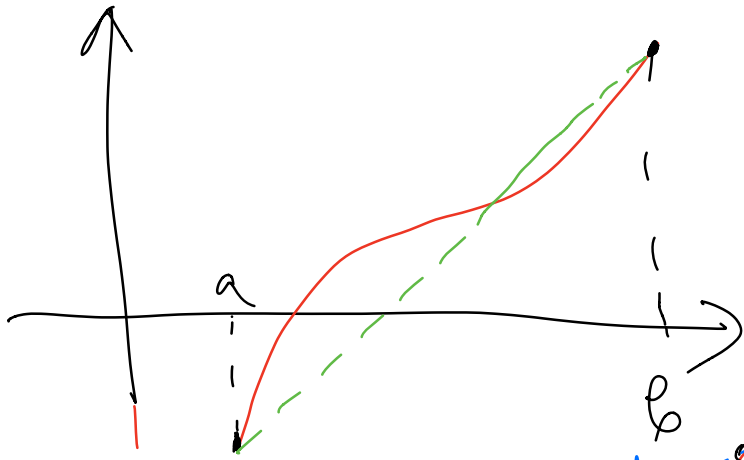
Solve  $p(x) = 0$  instead

Typically  $p(x)$  is linear  
or quadratic

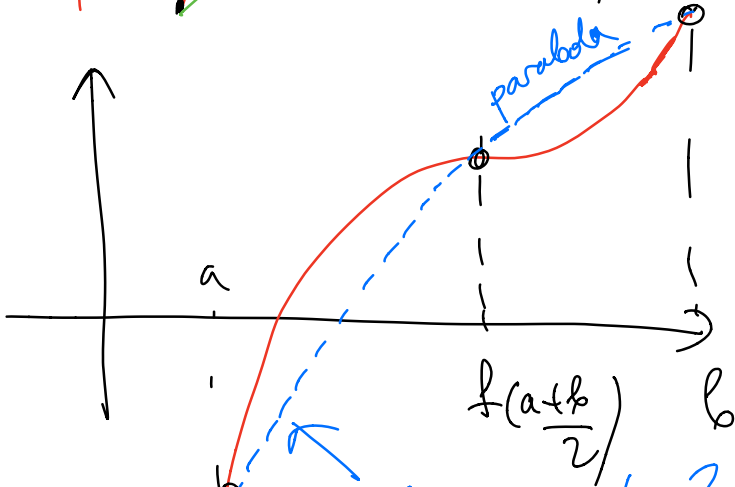
How do we find  $p(x)$ ?



Option 1

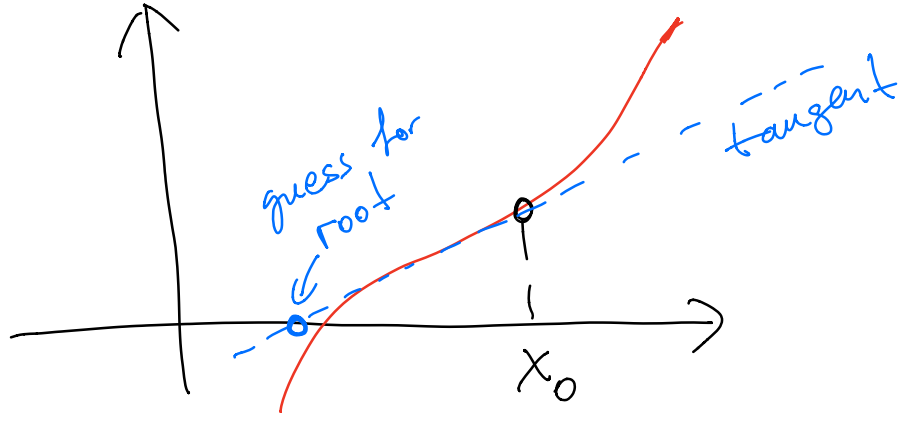


Interpolation



$P(x) = dx^2 + \beta x + \gamma$

Option 2 :



Option 2 : Taylor series

$$f(x) = \sum_{n=0}^{\infty} \frac{f^{(n)}(x_0)}{n!} (x-x_0)^n$$

$$p(x) = f(x_0) + f'(x_0)(x-x_0) + \frac{1}{2} f''(x_0)(x-x_0)^2 + \dots$$

$$p(x) = \sum_{n=0}^k \frac{f^{(n)}(x_0)}{n!} (x-x_0)^n$$

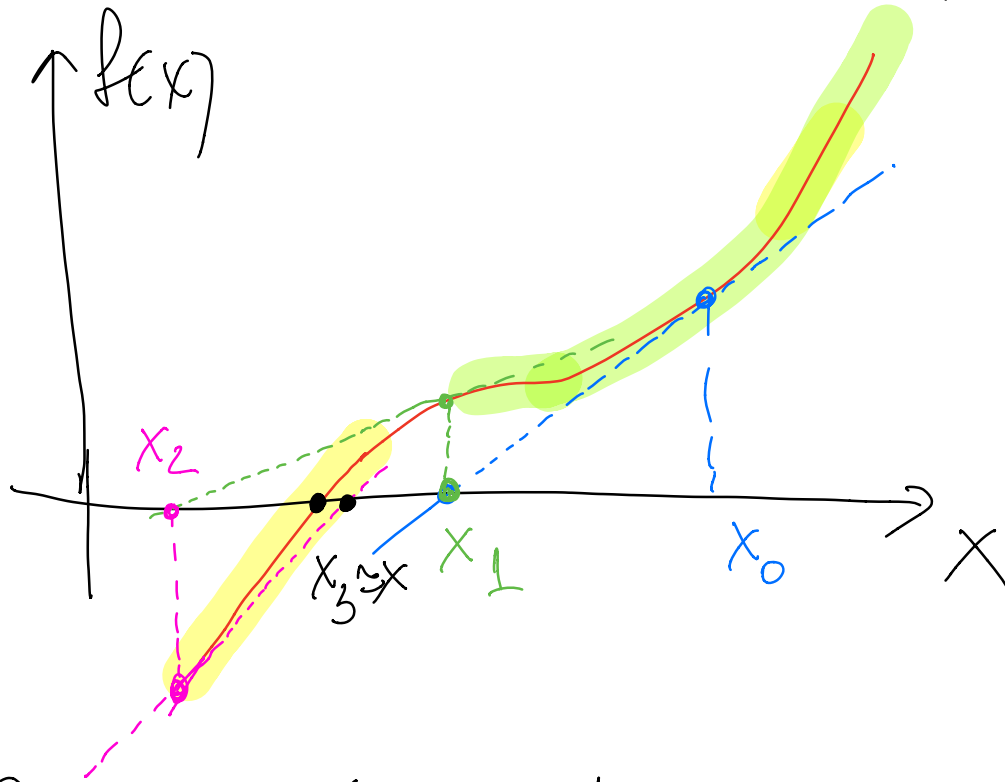
Truncated Taylor series

$$f(x) - p(x) = \frac{f^{(k+1)}(\xi)}{(k+1)!} (x-x_0)^{k+1}$$

Remainder

$\xi$  is between  $x$  and  $x_0$

# Newton's method



Given  $x_k, k = 0, 1, 2, \dots$   
Compute  $x_{k+1}$

$$f(x) \approx p_k(x) = f(x_k) + f'(x_k)(x - x_k)$$

$$p_k(x) = 0 \quad \text{solve for } x$$

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x_k)}$$

Newton's method

E.g.  $f(x) = x^2 - c$

$$f' = 2x$$

$$x_{k+1} = x_k - \frac{(x_k^2 - c)}{2x_k}$$

$$= \frac{1}{2} \left( x_k + \frac{c}{x_k} \right)$$

Babylonian  $(=)$  Newton's

Numerical  
Analysis: Include more  
terms in Taylor series

$$e_k = x_k - x$$

(sometimes  $e_k = |x - x_k|$ )

$$f(x) = 0 = f(x_k) + f'(x_k)(x - x_k) + \frac{1}{2}(x - x_k)^2 f''(\xi)$$

$\xi$  is between  $x$  and  $x_k$

$$x_k = e_k - x$$

Divide by  $f'(x_k)$

$$\left( x_k - \frac{f(x_k)}{f'(x_k)} \right) - x = \frac{1}{2} \frac{(x - x_k)^2 f''(\xi)}{f'(x_k)}$$

$$x_{k+1} - x = e_{k+1}$$

$$e_{k+1} = -\frac{1}{2} \frac{f''(\xi)}{f'(x_k)} e_k^2$$

$$\frac{e_{k+1}}{e_k^2} = -\frac{f''(\xi)}{2f'(x_k)} \rightarrow C$$

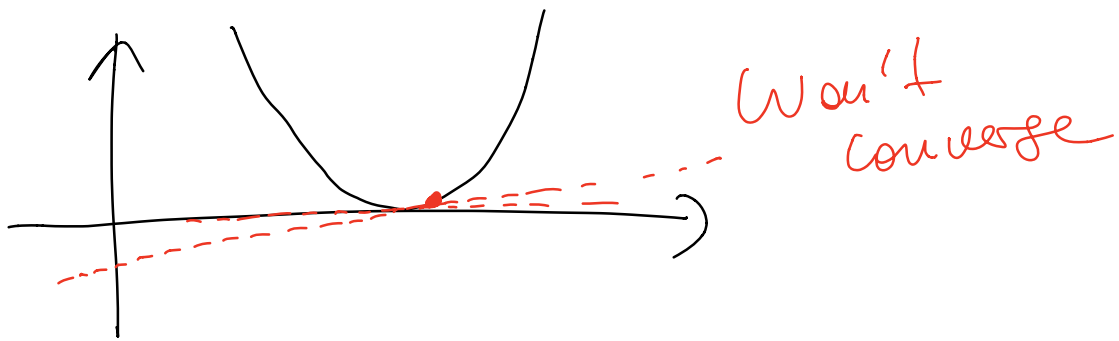
Assume method converges

$$x_k \rightarrow x$$

$$\xi \rightarrow x$$

$$\frac{|e_{k+1}|}{|e_k^2|} \rightarrow \frac{1}{2} \frac{f''(x)}{f'(x)} = C$$

Double root:  $f'(x) = 0$



$$|e_{k+1}| < |e_k|$$

We want  $|e_{k+1}| \leq \frac{|e_k|}{2}$

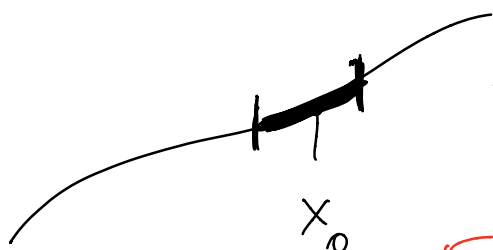
$$\left| \frac{f''(\xi)}{2f'(x_k)} \right| e_k^2 < \frac{|e_k|}{2}$$

$$|e_k| < \frac{|2f'(x_k)|}{|2f''(\xi)|} < A$$

$$e_0 = x_0 - x$$

$$|x_0 - x| < \left| \frac{2 f'(x_0)}{2 f''(\xi)} \right|$$

$\xrightarrow{\quad}$   
 $\xrightarrow{\quad}$



Theorem 1.8

Estimate :

Practice

$$|x_0 - x| < \left| \frac{f'(x)}{f''(x)} \right|$$

$A$

Assumption in Suli:

$$\left| \frac{f''(y_1)}{f'(y_2)} \right| < A$$

$$\forall y_1, y_2 \in [x - \delta, x + \delta]$$



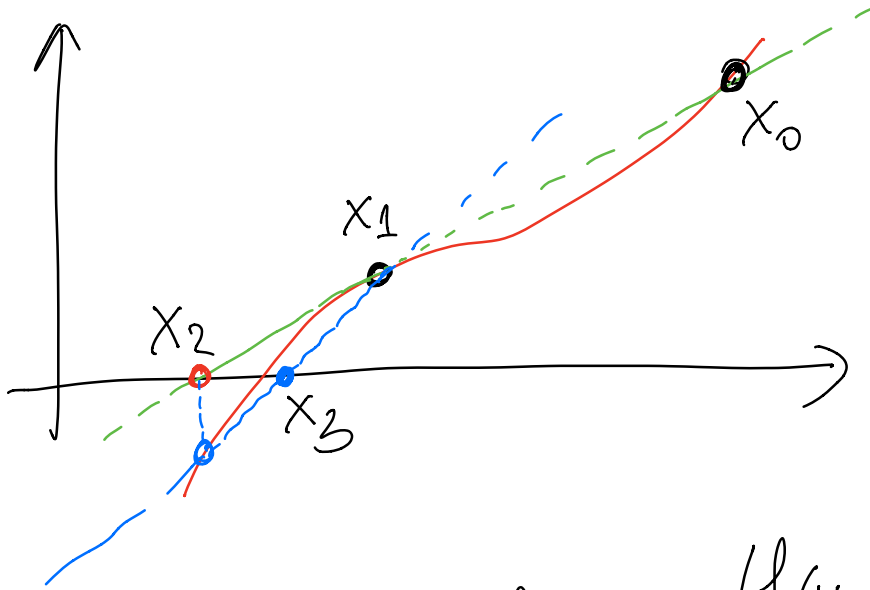
Theorem: If

$$|e_0| = |x_0 - x| < \frac{1}{A}$$

then Newton's method  
will converge

Newton's method needs to  
be started close to root  
and then it will converge  
fast

# Secant method



$$f(x) = p(x) = f(x_k) + \frac{(f(x_{k+1}) - f(x_k))}{(x_{k+1} - x_k)} (x - x_k)$$

$$\text{If } x = x_k \Rightarrow p(x_k) = f(x_k)$$

$$\text{If } x = x_{k+1} \Rightarrow p(x_{k+1}) = f(x_{k+1})$$

Newton  $p(x) = f(x_k) + f'(x_k)(x - x_k)$

Conclusion: As  $x_k \rightarrow x$

$$x_{k+1} \rightarrow x$$

These two become closer

$P(x) = 0$  solve for  $x_{k+2}$

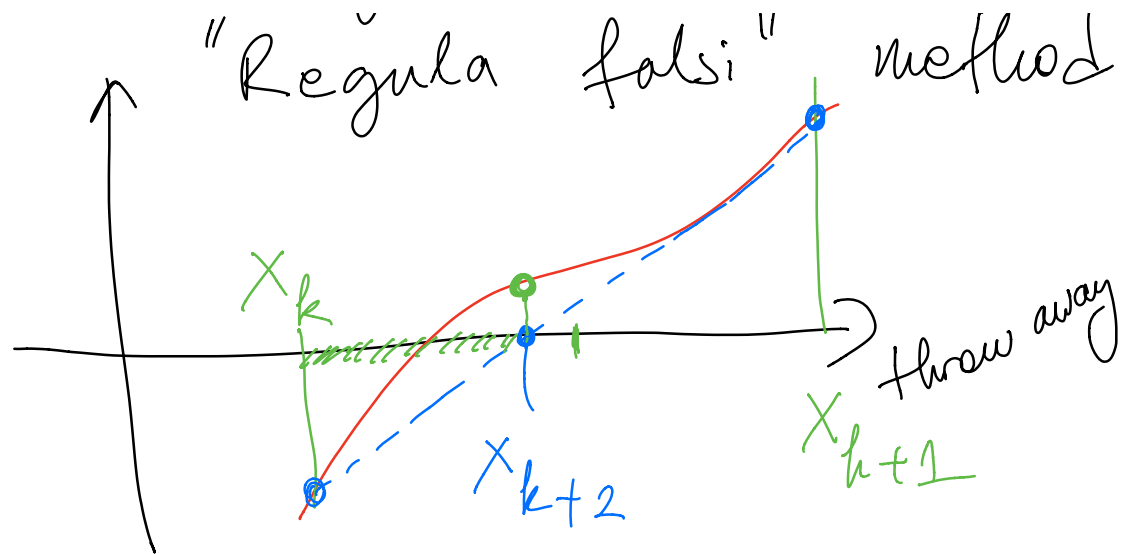
Secant method

$$x_{k+1} = x_k - \frac{(x_k - x_{k-1}) f_k}{(f_k - f_{k-1})}$$

$$f_k = f(x_k)$$

$$\frac{f(x)}{f'(x)}$$

We can combine  
Bisection + secant  
or "Safeguarded" secant



Problem 3 in HW 1:  
 Secant method:

$$|e_{k+1}| \rightarrow M |e_k|^q$$

Bisection:  $q = 1, \mu = 1/2$   
 } Steffensen  
 } Newton:  $q = 2, \mu = \frac{f''(x)}{2f'(x)}$   
 (quadratically convergent)  
 Secant:  $q = \frac{1}{2}(1 + \sqrt{5}) \approx 1.6$

# Fixed-point iteration

$$x_{k+1} = g(x_k)$$

$$x_k \rightarrow x, \quad f(x) = 0$$

$$x = g(x) \quad \begin{matrix} \nearrow \\ \text{equivalent} \end{matrix}$$

Babylonian

$$g(x) = \frac{1}{2} \left( x + \frac{c}{x} \right)$$

$$x = g(x) \Leftrightarrow x = \sqrt{c}$$

$$x = x^2 - c + x = g(x)$$

$$0 = x^2 - c = f(x)$$

$$x_{k+1} = g(x_k)$$

Fixed-point iteration

Example: Relaxation method

$$x_{k+1} = g(x_k) = x_k - \lambda f(x_k)$$

$$g(x) = x - \lambda f(x)$$

$\lambda \neq 0$

$$x = g(x) \Rightarrow$$

$$\cancel{x} = \cancel{x} - \lambda f(x) = 0$$

$$\Rightarrow f(x) = 0$$

Method is consistent

What is a good  $\lambda$

$$x_{k+1} = x_k - \lambda f(x_k)$$

---


$$x_{k+1} = g(x_k) \leftarrow \begin{array}{l} \text{Taylor} \\ \text{Series around} \\ \underline{x} \text{ (solution)} \end{array}$$

$$= g(x) + (x_k - x)g'(\xi)$$

$\xi$  is in-between  $x, x_k$

Express in terms of error

$$\left\{ \begin{array}{l} x_{k+1} = x + \underset{\substack{\uparrow \\ \text{error}}}{e_{k+1}} \\ x_k = x + e_k \end{array} \right.$$

$$e_{k+1} = x_{k+1} - x =$$


---

$$= \overbrace{g(x)}^x + (x_k - x) g'(\xi)$$

$$[\text{we know } x = g(x)] \Leftrightarrow f(x) = 0$$

$$= \underbrace{(x_k - x)}_{e_k} g'(\xi)$$

$$= e_k g'(\xi)$$

$$\underbrace{|e_{k+1}|} = \underbrace{|g'(\xi)|} \underbrace{|e_k|}$$

We want

$$|e_{k+1}| < |e_k|$$

$$\Rightarrow |g'(\xi)| < 1 \quad (*)$$



$\xi$  is between  $x$  and  $x_k$

(\*) should be true if  $\xi$  is close to root

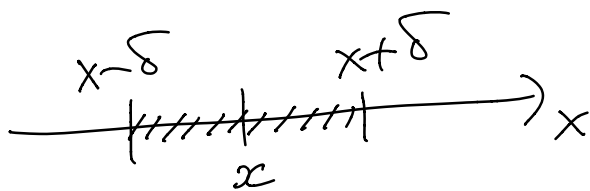
Theorem:

$\exists f \quad g \in C^1 \leftarrow$  continuously differentiable

$|g'(y)| < 1 \quad \forall y \in [x-\delta, x+\delta]$

where

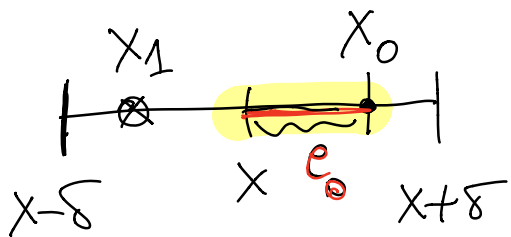
$$x = g(x)$$



then fixed-point iteration

$$x_{h+1} = g(x_h)$$

will converge to  $x$  if  
"Proof":  $x_0 \in [x-\delta, x+\delta]$



$$|e_0| < \delta$$

$$|e_1| = |g'(\xi_0)| |e_0|$$

$$|e_1| < |e_0| < \delta$$

$$x_1 \in [x-\delta, x+\delta]$$

$$\Rightarrow |e_2| = |g'(\xi_1)| |e_1|$$

$$|e_2| < |e_1| < |e_0|$$

By induction

$$|e_{k+1}| \leq L |e_k|$$

$$0 < L < 1$$

$$L = \max_{y \in [x-\delta, x+\delta]} |g'(y)| < 1$$

$$|e_k| \leq L^k |e_0|$$
$$L^k \rightarrow 0$$

$$\Rightarrow |e_k| \rightarrow 0 \quad \text{Q.E.D.}$$

$$\frac{|e_{k+1}|}{|e_k|} \rightarrow |g'(x)|$$

since  $x_k \rightarrow x$   
 $\xi \rightarrow x$

Fast convergence we want

$$|g'(x)| \ll 1$$

Ideally

$$g'(x) = 0$$

convergence is faster than

linear  $\Rightarrow$  super linear

convergence.

Example : Relaxation

$$x_{k+1} = x_k - \lambda f(x_k)$$

$$g(x) = x - \lambda f(x)$$

$$g'(x) = 1 - \lambda f'(x) = 0$$

$$\lambda = -\frac{1}{f'(x)}$$

Best  $\lambda$ . Fastest  
(super linear) convergence for

Doesn't work

$$x_{k+1} = x_k - \frac{f(x_k)}{f'(x)}$$

we don't know  $x$

c.f. Newton

$$x_k - \frac{f(x_k)}{f'(x_k)}$$

$\rightarrow$

For Newton's method

$$g(x) = x - \frac{f(x)}{f'(x)}$$

What is  $g'(x)$  at the root  $x$  that satisfies  $x=g(x)$

$$\begin{aligned} g'(x) &= \cancel{x} - \frac{\cancel{f'(x)}}{\cancel{f'(x)}} + \frac{f(x) f''(x)}{(f'(x))^2} \\ &= \frac{f(x) f''(x)}{(f'(x))^2} \end{aligned}$$

$$\text{If } f(x) = 0 \Rightarrow g'(x) = 0$$

Consistent with Newton's method converging quadratically.